

# MATLAB Examples: Linear Block Codes

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## 1 The Galois Field $\mathbb{F}_p$ for Prime $p$

Currently this document just gives snippets of example code which should help you get started. If you need help, keep the “help” and “lookfor” commands in mind.

### 1.1 A Few Commands

```
>> help eye
```

EYE Identity matrix.

EYE(N) is the N-by-N identity matrix.

EYE(M,N) or EYE([M,N]) is an M-by-N matrix with 1's on the diagonal and zeros elsewhere.

EYE(SIZE(A)) is the same size as A.

EYE with no arguments is the scalar 1.

EYE(M,N,CLASSNAME) or EYE([M,N],CLASSNAME) is an M-by-N matrix with 1's of class CLASSNAME on the diagonal and zeros elsewhere.

Example:

```
x = eye(2,3,'int8');
```

See also SPEYE, ONES, ZEROS, RAND, RANDN.

```
>> help mod
```

MOD Modulus after division.

MOD(x,y) is  $x - n \cdot y$  where  $n = \text{floor}(x./y)$  if  $y \neq 0$ . If  $y$  is not an integer and the quotient  $x./y$  is within roundoff error of an integer, then  $n$  is that integer. The inputs  $x$  and  $y$  must be real arrays of the same size, or real scalars.

The statement “ $x$  and  $y$  are congruent mod  $m$ ” means  $\text{mod}(x,m) == \text{mod}(y,m)$ .

By convention:

MOD(x,0) is  $x$ .

MOD(x,x) is 0.

MOD(x,y), for  $x \neq y$  and  $y \neq 0$ , has the same sign as  $y$ .

Note: REM(x,y), for  $x \neq y$  and  $y \neq 0$ , has the same sign as  $x$ .

MOD(x,y) and REM(x,y) are equal if  $x$  and  $y$  have the same sign, but differ by  $y$  if  $x$  and  $y$  have different signs.

See also REM.

Overloaded functions or methods (ones with the same name in other directories)  
help sym/mod.m

>> help de2bi

DE2BI Convert decimal numbers to binary numbers.

B = DE2BI(D) converts a nonnegative integer decimal vector D to a binary matrix B. Each row of the binary matrix B corresponds to one element of D. The default orientation of the of the binary output is Right-MSB; the first element in B represents the lowest bit.

In addition to the vector input, three optional parameters can be given:

B = DE2BI(...,N) uses N to define how many digits (columns) are output.

B = DE2BI(...,N,P) uses P to define which base to convert the decimal elements to.

B = DE2BI(...,FLAG) uses FLAG to determine the output orientation. FLAG has two possible values, 'right-msb' and 'left-msb'. Giving a 'right-msb' FLAG does not change the function's default behavior. Giving a 'left-msb' FLAG flips the output orientation to display the MSB to the left.

Examples:

D = [12; 5];

B = de2bi(D)

B =  
0 0 1 1  
1 0 1 0

B = de2bi(D,5)

B =  
0 0 1 1 0  
1 0 1 0 0

T = de2bi(D,[],3)

T =  
0 1 1  
2 1 0

B = de2bi(D,5,'left-msb')

B =  
0 1 1 0 0  
0 0 1 0 1

See also BI2DE.

>> help dec2base

DEC2BASE Convert decimal integer to base B string.

DEC2BASE(D,B) returns the representation of D as a string in base B. D must be a non-negative integer array smaller than  $2^{52}$  and B must be an integer between 2 and 36.

DEC2BASE(D,B,N) produces a representation with at least N digits.

Examples

dec2base(23,3) returns '212'

dec2base(23,3,5) returns '00212'

See also BASE2DEC, DEC2HEX, DEC2BIN.

>> help nchoosek

NCHOOSEK Binomial coefficient or all combinations.

NCHOOSEK(N,K) where N and K are non-negative integers returns  $N!/K!(N-K)!$ . This is the number of combinations of N things taken K at a time.

When a coefficient is greater than  $10^{15}$ , a warning will be produced indicating possible inexact results. In such cases, the result is good

to 15 digits.

NCHOOSEK(V,K) where V is a vector of length N, produces a matrix with  $N!/K!(N-K)!$  rows and K columns. Each row of the result has K of the elements in the vector V. This syntax is only practical for situations where N is less than about 15.

Class support for inputs N,K,V:  
float: double, single

See also PERMS.

## 1.2 Now For Some Coding

```
>> n = 6;  
>> k = 3;  
>> p = 2;  
>> In = eye(n);  
>> Ik = eye(k);  
>> Ink = eye(n-k);
```

```
>> P = [1 1 0;0 1 1;1 0 1]
```

P =

```
    1    1    0  
    0    1    1  
    1    0    1
```

```
>> G = [Ik P];  
>> H = mod([-P' Ink],p)
```

H =

```
    1    0    1    1    0    0  
    1    1    0    0    1    0  
    0    1    1    0    0    1
```

```
>> mod(G*H',p) % Test G and H construction
```

ans =

```
    0    0    0  
    0    0    0  
    0    0    0
```

## 1.3 Encoding and Listing Codewords

We note that “u = dec2base(0:(p<sup>k</sup> - 1),p,k)-'0'” can be used instead of “de2bi” for  $p > 2$ .

```
>> u = de2bi(0:(2k - 1),k) % List all binary input vectors
```

u =

```
    0    0    0  
    1    0    0
```

```

0    1    0
1    1    0
0    0    1
1    0    1
0    1    1
1    1    1

```

```
>> C = mod(u*G,p) % List all codewords
```

```
C =
```

```

0    0    0    0    0    0
1    0    0    1    1    0
0    1    0    0    1    1
1    1    0    1    0    1
0    0    1    1    0    1
1    0    1    0    1    1
0    1    1    1    1    0
1    1    1    0    0    0

```

## 1.4 Syndromes

```
>> N2 = nchoosek(1:n,2)
```

```
ans =
```

```

1    2
1    3
1    4
1    5
1    6
2    3
2    4
2    5
2    6
3    4
3    5
3    6
4    5
4    6
5    6

```

```
>> E2 = zeros(length(N2),n);
```

```
>> for i=1:length(N2); E2(i,N2(i,:)) = 1; end % All weight 2 error patterns
```

```
>> E2
```

```
E2 =
```

```

1    1    0    0    0    0
1    0    1    0    0    0
1    0    0    1    0    0
1    0    0    0    1    0
1    0    0    0    0    1
0    1    1    0    0    0
0    1    0    1    0    0

```

```

0 1 0 0 1 0
0 1 0 0 0 1
0 0 1 1 0 0
0 0 1 0 1 0
0 0 1 0 0 1
0 0 0 1 1 0
0 0 0 1 0 1
0 0 0 0 1 1

```

```
>> S2 = mod(E2*H',2) % List syndromes of weight 2 patterns
```

```
S2 =
```

```

1 0 1
0 1 1
0 1 0
1 0 0
1 1 1
1 1 0
1 1 1
0 0 1
0 1 0
0 0 1
1 1 1
1 0 0
1 1 0
1 0 1
0 1 1

```

```
>> S2int = bi2de(S2); % Assign each syndrome to an integer between 0 and 2^(n-k) - 1
```

## 1.5 Simulation

```
>> M = 5; % Handle M transmissions at once
```

```
>> msg = floor(rand(M,1)*2^k) % Generate uniform random message numbers
```

```
msg =
```

```

4
6
7
5
1

```

```
>> u = de2bi(msg,k) % Map message number to bit vector
```

```
u =
```

```

0 0 1
0 1 1
1 1 1
1 0 1
1 0 0

```

```
>> c = mod(u*G,p); % Encode each message
```

```
>> noise = rand(M,n)<0.1 % Generate BSC noise with error prob. 0.1
```

```
noise =
    0    0    0    0    1    0
    0    1    0    0    0    0
    0    0    0    0    0    0
    0    0    0    0    1    0
    1    0    1    0    0    0
```

```
>> recv = mod(c+noise,p);
```

## 1.6 Matrix Tricks

The following tricks MATLAB into performing matrix inverses over prime fields. It uses the fact that  $\det(A)A^{-1}$  is an integer matrix if  $A$  is an integer matrix and it uses the fact that  $a^{p-2} = a^{-1}$  for all  $a \in GF(p)$ . Due to the finite precision of IEEE doubles, the first trick may fail if any element of  $\det(A)A^{-1}$  is greater than  $10^{16}$ . Likewise, the second may fail if  $(p-1)^{p-2} > 10^{16}$ .

```
>> A = floor(2*rand(5,5)) % Generate random 5 by 5 binary matrix
```

```
A =
    1    0    1    0    1
    0    0    1    0    1
    0    1    0    0    1
    0    0    0    1    0
    0    0    0    1    1
```

```
>> det(A)
```

```
ans =
```

```
-1
```

```
>> invA = mod(inv(A)*det(A),2) % Modulo 2 inverse trick (det(A) must be odd)
```

```
invA =
```

```
    1    1    0    0    0
    0    0    1    1    1
    0    1    0    1    1
    0    0    0    1    0
    0    0    0    1    1
```

```
>> mod(invA*A,2) % Verify that it works
```

```
ans =
```

```
    1    0    0    0    0
    0    1    0    0    0
    0    0    1    0    0
    0    0    0    1    0
    0    0    0    0    1
```

```
>> A = floor(7*rand(5,5)) % Generate random 5 by 5 matrix over GF(7)
```

```
A =
```

```

5   6   1   2   6
1   6   0   1   4
6   2   3   6   1
0   3   5   4   4
3   3   2   1   4

```

```
>> det(A)/7 % Check determinant not divisible by 7
```

```
ans =
```

```
66.4286
```

```
>> invA = round(mod(mod(round(inv(A)*det(A)),7)*mod(det(A),7)^5,7))
```

```
invA =
```

```

6   0   4   3   1
5   6   0   5   6
5   5   3   1   5
1   2   1   0   5
3   3   4   4   0

```

```
>> mod(invA*A,7) % Test inverse
```

```
ans =
```

```

1   0   0   0   0
0   1   0   0   0
0   0   1   0   0
0   0   0   1   0
0   0   0   0   1

```

## 2 Extension Fields $\mathbb{F}_{2^m}$

### 2.1 A Few Commands

Matlab has built in routines that work for extension fields of characteristic 2. These commands can be listed by typing "help gfhelp".

**GF** Create a Galois field array.

`X_GF = GF(X,M)` creates a Galois field array from `X` in the field  $\text{GF}(2^M)$ , for  $1 \leq M \leq 16$ . The elements of `X` must be integers between 0 and  $2^M - 1$ . `X_GF` behaves like a MATLAB array, and you can use standard indexing and arithmetic operations (+, \*, .\*, .^, \, etc.) on it. For a complete list of operations you can perform on `X_GF`, type "GFHELP".

`X_GF = GF(X,M,PRIM_POLY)` creates a Galois field array from `X` and uses the primitive polynomial `PRIM_POLY` to define the field. `PRIM_POLY` must be a primitive polynomial in decimal representation. For example, the polynomial  $D^3 + D^2 + 1$  is represented by the number 13, because 1 1 0 1 is the binary form of 13.

`X_GF = GF(X)` uses a default value of  $M = 1$ .

Example:

```
A = gf(randint(4,4,8,873),3); % 4x4 matrix in GF(2^3)
B = gf(1:4,3)'; % A 4x1 vector
C = A*B
```

C = GF(2^3) array. Primitive polynomial = 1+D+D^3 (11 decimal)

Array elements =

```
3
3
6
7
```

See also GFHELP, GFTABLE.

## 2.2 Standard Codes

```
>> n = 5;
>> k = 3;
>> m = 2;
>> In = gf(eye(n),m);
>> Ik = gf(eye(k),m);
>> Ink = gf(eye(n-k),m);

>> P = gf([1 1;1 2;1 3],m) % (5,3) Hamming code over GF(4)
```

P = GF(2^2) array. Primitive polynomial = D^2+D+1 (7 decimal)

Array elements =

```
1 1
1 2
1 3
```

```
>> G = [Ik P];
>> H = [P' Ink];
```

H = GF(2^2) array. Primitive polynomial = D^2+D+1 (7 decimal)

Array elements =

```
1 1 1 1 0
1 2 3 0 1
```

```
>> G*H' % Test G and H construction
```

ans = GF(2^2) array. Primitive polynomial = D^2+D+1 (7 decimal)

Array elements =

```
0 0
0 0
0 0
```



## 2.3 Reed-Solomon Codes via FFTs

```
% Reed-Solomon code over GF(256)
m = 8;
q = 2^m;
n = q-1;
r = 6;
k = n-r;

% Encode message using FFT
u = gf([floor(rand(1,k)*q) zeros(1,n-k)],m);
x = fft(u);

% Construct random error pattern of weight "ne"
ne = 4;
e = gf(zeros(1,n),m);
loc = randperm(n);
mag = gf(floor(rand(1,ne)*(q-1))+1,m);
e(loc(1:ne)) = mag;

% Add errors and compute syndrome via IFFT
y = x+e;
syn = ifft(y);
syn = syn((k+1):n);
```